

# Numerical Optimization Methods

## 5. Least Squares Methods

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# 1. Preliminary Concepts

- Linear system vs. matrix form
- Backward substitution of solving  $Ax = b$

– Given a upper triangle  $A = \begin{pmatrix} a_{11} & \cdots & a_{1n} \\ & \ddots & \vdots \\ & & a_{nn} \end{pmatrix}$

$$x_k = \frac{b_k - \sum_{j=k+1}^n a_{kj}x_j}{a_{kk}}, k = n, \dots, 1$$

– For a lower triangle, forward substitution!

- Gaussian elimination

- Augmented matrix,  $A' = (A|b)$

$$A'(2,:) = A'(2,:) - \frac{a_{21}}{a_{11}} A'(1,:)$$

$$A'(k,:) = A'(k,:) - \frac{a_{k1}}{a_{11}} A'(1,:)$$

$$k = 2, \dots, n$$

$$A'(3,2:n) = A'(3,2:n) - \frac{a_{32}}{a_{22}} A'(3,2:n)$$

$$A'(k, 2:n) = A'(k, 2:n) - \frac{a_{k2}}{a_{22}} A'(k, 2:n)$$

$$k = 3, \dots, n$$

Convert A into an Upper triangle, .....

- **Assumption**: the diagonal element  $a_{ii}$  must always be away from Zero!

- LU decomposition (square matrix)
  - Recall Gaussian elimination procedure

Eliminating each column (i.e. zero out each column),

$$M^{(1)} = \begin{pmatrix} 1 & & & \\ -l_{21} & 1 & & \\ \vdots & & \ddots & \\ -l_{n1} & 0 & \dots & 1 \end{pmatrix},$$

$$l_{21} = \frac{a_{21}}{a_{11}}, \dots, l_{n1} = \frac{a_{n1}}{a_{11}}$$

..... $M^{(i)}$ .....

$$U = M^{(n-1)} \dots M^{(1)} A$$

$$L = M^{(1)-1} \dots M^{(n-1)-1}$$

$$A = LU$$

**Assumption:** diagonal elements of A away from zero!

– Apply LU factorization to  $Ax = b$

$$A = LU,$$

$Ly = b$ , forward substitution

$Ux = y$ , backward substitution

– Compare with Gaussian Elimination,

$A$  decomposition is independent of  $b$ !

– Problem

$$A = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, \quad \det(A) = -1$$

Or,  $a_{ii}$  is near zero!

- Pivoting for zero out diagonal elements
  - Partial pivoting,

Choosing  $q$ , such that,

$$|a_{qk}^{(k-1)}| = \max_{k \leq i \leq n} |a_{ik}^{(k-1)}|$$

$$\begin{array}{l}
 k \rightarrow \\
 q \rightarrow
 \end{array}
 \begin{pmatrix}
 a_{11} & & \dots & & a_{1n} \\
 0 & \ddots & & & \\
 \vdots & \ddots & a_{kk} & & \vdots \\
 & 0 & \vdots & \ddots & \\
 0 & \dots & a_{nk} & \dots & a_{nn}
 \end{pmatrix}$$

Interchanging rows  $k$  and  $q$  forms a permutation!

Cannot guarantee stability of  $Ax = b$ !

– Complete pivoting,

$$|a_{qr}^{(k-1)}| = \max_{k \leq i, j \leq n} |a_{ij}^{(k-1)}|$$

$$\begin{array}{l}
 k \rightarrow \\
 q \rightarrow
 \end{array}
 \begin{pmatrix}
 a_{11} & \dots & & & a_{1n} \\
 0 & \ddots & & & \\
 \vdots & \ddots & a_{kk} & & \vdots \\
 & 0 & \vdots & \ddots & \\
 0 & \dots & a_{nk} & \dots & a_{nn}
 \end{pmatrix}$$

$k \quad r$

Interchanging rows  $k$  and  $q$  and then columns  $k$  and  $r$ !

– Very expensive! Have to pay attention to special matrices!

- Special matrices (no interchanging!)
  - Symmetric Positive Definite matrix,  $x^T Ax > 0, x \neq 0$

$$A = LU = L \text{diag}(u_{ii}) \begin{pmatrix} 1 & \frac{u_{12}}{u_{11}} & \dots & \frac{u_{1n}}{u_{11}} \\ & 1 & & \frac{u_{2n}}{u_{22}} \\ & & \ddots & \vdots \\ & & & 1 \end{pmatrix}$$

$$L = \begin{pmatrix} 1 & & & \\ \frac{u_{12}}{u_{11}} & 1 & & \\ \vdots & & \ddots & \vdots \\ \frac{u_{1n}}{u_{11}} & \frac{u_{2n}}{u_{22}} & \dots & 1 \end{pmatrix}$$

$$A = L \text{diag}(u_{ii}) L^T$$



– Cholesky factorization

$$G = L\sqrt{\text{diag}(a_{ii})}$$
$$A = GG^T$$

– SPD matrix's properties,

- Invertible;
- $a_{ii} > 0$ ;
- $\max_{1 \leq k, j \leq n} |a_{kj}| \leq \max_{1 \leq i \leq n} |a_{ii}|$ ;
- $(a_{ij})^2 < a_{ii}a_{jj}, i \neq j$ ;

– Diagonally dominate matrix,

$$a_{ii} \geq \sum_{j=1, j \neq i}^n |a_{ij}|$$

## – Banded matrix

Tridiagonal matrix,

$$A = \begin{pmatrix} a_{11} & a_{12} & 0 & \dots & 0 \\ a_{21} & a_{22} & a_{23} & \ddots & \vdots \\ 0 & a_{32} & a_{33} & \ddots & 0 \\ \vdots & \ddots & \ddots & \ddots & a_{n-1,n} \\ 0 & \dots & 0 & a_{n,n-1} & a_{nn} \end{pmatrix}$$

If  $|a_{11}| > |a_{12}|$ ,  $|a_{ii}| \geq |a_{i,i-1}| + |a_{i,i+1}|, \dots |a_{nn}| > |a_{n,n-1}|$ ,

apply LU to  $A$  directly!

- QR Decomposition

Let  $A \in R^{m \times n}$ ,  $m \geq n$ , be full rank,

$$A = (Q_1 \quad Q_2)_{m \times m} \begin{pmatrix} T_{upper} \\ 0 \end{pmatrix},$$

$Q_1 \in R^{m \times n}$ ,  $Q_2 \in R^{m \times (m-n)}$ ,  $T_{upper} \in R^{n \times n}$ ,

$$Q_1^T Q_1 = I, Q_2^T Q_2 = I, Q_1^T Q_2 = 0$$

This is the column orthogonal!

– How to compute QR?

Givens rotation,

$$G(i, j, \theta) = \begin{bmatrix} 1 & \dots & 0 & \dots & 0 & \dots & 0 \\ \vdots & \ddots & \vdots & & \vdots & & \vdots \\ 0 & \dots & c & \dots & -s & \dots & 0 \\ \vdots & & \vdots & \ddots & \vdots & & \vdots \\ 0 & \dots & s & \dots & c & \dots & 0 \\ \vdots & & \vdots & & \vdots & \ddots & \vdots \\ 0 & \dots & 0 & \dots & 0 & \dots & 1 \end{bmatrix},$$

Applying a series of Givens rotations to  $A$  leads to a upper matrix  $T_{upper}$ !

–Example:

For an underdetermined linear system,  $Ax = b$ ,  $A \in R^{m \times n}$ ,  $m < n$ ,

$$A^T = (Q_1 \quad Q_2) \begin{pmatrix} T \\ 0 \end{pmatrix}$$
$$(T^T, 0)(Q_1 \quad Q_2)^T x = b$$
$$x' = (Q_1 \quad Q_2) \begin{pmatrix} T^{-T} \\ 0 \end{pmatrix} b = Q_1 T^{-T} b$$
$$Ax' = b$$

For nullspace,  $Ax = 0$ ,  $Q_2$  forms a basis, the general solution is,

$$x = x' + Q_2 y, \quad \forall y \in R^{n-m}$$

- SVD and Generalized SVD

- SVD

For  $A \in R^{n \times m}$ , there exist orthogonal matrices,

$U \in R^{n \times n}$  and  $V \in R^{m \times m}$ ,

$U^T A V = \Sigma = \text{diag}(\sigma_1, \dots, \sigma_p)$ ,

- Properties,

(1)  $A v_i = \sigma_i u_i, \quad i = 1 \dots p$ ,

(2)  $\text{Rank}(A) = p, p = \min(n, m)$

$\text{Null}(A) = \text{span}\{v_{p+1}, \dots, v_m\}$

$\text{Range}(A) = \text{span}\{u_1, \dots, u_p\}$

$$A = U \Sigma V^*$$

$$= \begin{bmatrix} \mathbf{U}_{\mathcal{R}} & \mathbf{U}_{\mathcal{N}} \end{bmatrix} \begin{bmatrix} \sigma_1 & 0 & \dots & & \dots & 0 \\ 0 & \sigma_2 & & & & \\ \vdots & & \ddots & & & \\ & & & \sigma_p & & \\ \hline & & & & 0 & \\ \vdots & & & & & \ddots \\ 0 & & & & & 0 \end{bmatrix} \begin{bmatrix} \mathbf{V}_{\mathcal{R}}^* \\ \mathbf{V}_{\mathcal{N}}^* \end{bmatrix} = \begin{bmatrix} u_1 & \dots & u_p & u_{p+1} & \dots & u_m \end{bmatrix} \begin{bmatrix} \mathbf{S}_{\rho \times \rho} & \mathbf{0} \\ \mathbf{0} & \mathbf{0} \end{bmatrix} \begin{bmatrix} v_1^* \\ \vdots \\ v_p^* \\ v_{p+1}^* \\ \vdots \\ v_n^* \end{bmatrix}$$

(3) For two  $A, B \in R^{n \times m}$ ,

$$\sigma_{\max}(A + B) \leq \sigma_{\max}(A) + \|B\|_2$$

$$\sigma_{\min}(A + B) \geq \sigma_{\min}(A) - \|B\|_2$$

## – Generalized SVD

○ For  $A \in R^{n_1 \times m}, B \in R^{n_2 \times m}, n_1 > m$ , assume that

$$\text{rank} \begin{pmatrix} A \\ B \end{pmatrix} = r$$

There exist orthogonal,  $U_1 \in R^{n_1 \times n_1}, U_2 \in R^{n_2 \times n_2}$ , and invertible  $X \in R^{m \times m}$ , such that,

$$U_1^T A X = D_A = \begin{pmatrix} I & & \\ & \text{diag}(\alpha_{p+1}, \dots, \alpha_r) & \\ & & 0 \end{pmatrix}$$
$$U_2^T B X = D_B = \begin{pmatrix} 0 & & \\ & \text{diag}(\beta_{p+1}, \dots, \beta_r) & \\ & & 0 \end{pmatrix}$$

where  $p = \max\{r - n_2, 0\}$ .

- Numerical algebra

- Orthogonal matrices for  $Ax = b$

$$x = A^{-1}b = A^T b, \quad O(2n^2)$$

If  $A = I - 2uu^T, u \in R^n,$

$$x = A^{-1}b = b - 2(u^T b)u, \quad O(4n)$$

- Permutation matrices  $Ax = b$

$$A_{ij} = \begin{cases} 1, & j = \pi_i \\ 0, & \text{otherwise} \end{cases}$$

$\pi = (\pi_1, \dots, \pi_n)$  is a permutation of  $(1, 2, \dots, n)$ .

○ Example,

There are six permutations of the set  $\{1,2,3\}$ , namely  $(1,2,3), (1,3,2), (2,1,3), (2,3,1), (3,1,2)$ , and  $(3,2,1)$ !

Each  $\pi_i$  contains only one entry with value 1 while others with 0;

$$Ax = (x_{\pi_1}, \dots, x_{\pi_n})$$

$A$  is orthogonal!

$$x = A^T b$$

Permuting the entries of  $b$ !



– Factor solver

Let  $A = A_1 \dots A_k$ , for  $Ax = b$ ,

$$y_1 := A_1^{-1}b$$

$$y_2 := A_2^{-1}y_1 \dots \dots$$

$$x = A_k^{-1}y_{k-1}$$

– Block  $LDL^T$  factorization

$A = PLDL^T P^T$ ,  $D$  is Block diagonal

○ Solve  $Ax = b$  by Gaussian elimination,

solve  $Py_1 = b$ ;

solve  $Ly_2 = y_1$ , forward substitution;

solve  $Dy_3 = y_2$ ;

solve  $L^T y_4 = y_3$ , backward substitution;

solve  $P^T x = y_4$ ;

## – Schur complement

Let  $X = \begin{bmatrix} A & B \\ B^T & C \end{bmatrix}_{n \times n}$ ,  $A \in R^{k \times k}$  symmetric and nonsingular, Schur complement of  $A$  in  $X$  is,  
$$S = C - B^T A^{-1} B$$

If  $C$  is invertible, Schur complement of  $C$  in  $X$  is,  
$$S = A - B^T C^{-1} B$$

○ Inverse of  $X$ ,

$$\begin{bmatrix} A & B \\ B^T & C \end{bmatrix}^{-1} = \begin{bmatrix} A^{-1} + A^{-1} B S^{-1} B^T A^{-1} & -A^{-1} B S^{-1} \\ -S^{-1} B^T A^{-1} & S^{-1} \end{bmatrix}$$

Positive definite  $X \leftrightarrow A$  and  $S$  are positive definite;

If  $A$  is positive definite, then  $X$  is positive semi-definite  $\leftrightarrow S$  is positive semi-definite!

## – Block elimination

To solve  $Ax = b$ , let  $A = \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix}$ ,  $x = \begin{bmatrix} x_1 \\ x_2 \end{bmatrix}$ ,  $b = \begin{bmatrix} b_1 \\ b_2 \end{bmatrix}$ , then

$$\begin{cases} x_1 = A_{11}^{-1}b_1 - A_{11}^{-1}A_{12}x_2 \\ (A_{22} - A_{21}A_{11}^{-1}A_{12})x_2 = b_2 - A_{21}A_{11}^{-1}b_1 \end{cases}$$

○ Schur complement of 1<sup>st</sup> block  $A_{11}$ ,

$$S = (A_{22} - A_{21}A_{11}^{-1}A_{12})$$

Construct sub terms,  $A_{11}^{-1}A_{12}$  and  $A_{11}^{-1}b_1$ ;

Form  $S$  and  $b' = b_2 - A_{21}A_{11}^{-1}b_1$ ;

Solve  $Sx_2 = b'$ ;

Solve  $A_{11}x_1 = b_1 - A_{12}x_2$ ;

- Solving underdetermined linear system
  - Given,  $Ax = b$ ,  $A \in R^{m \times n}$ ,  $m < n$ ,  $\text{rank}(A) = m$
  - Solution form,
 
$$\{Fz + \hat{x} \mid z \in R^{n-m}\}$$
 where  $F$  is a matrix whose columns form a basis for nullspace of  $A$ , and  $\hat{x}$  is a particular solution.
  - How to construct  $F$  and  $\hat{x}$ ? (**splitting method**)
  - Partition  $A$ ,

$$Ax = (A_1, A_2) \begin{pmatrix} x_1 \\ x_2 \end{pmatrix} = b$$

where  $A_1 \in R^{m \times m}$  is nonsingular.

$$x_1 = A_1^{-1}b - A_1^{-1}A_2x_2$$

– General solution,

$$x = \begin{pmatrix} -A_1^{-1}A_2 \\ I \end{pmatrix} x_2 + \begin{pmatrix} A_1^{-1}b \\ 0 \end{pmatrix}$$

$$\text{then } F = \begin{pmatrix} -A_1^{-1}A_2 \\ I \end{pmatrix}, \hat{x} = \begin{pmatrix} A_1^{-1}b \\ 0 \end{pmatrix}!$$

Challenge is HOW to partition  $A$ ! (partition examples)

○ Apply QR factorization,  $A^T = (Q_1, Q_2) \begin{pmatrix} B \\ 0 \end{pmatrix}$

Substituting factors,

$$\hat{x} = Q_1 B^{-T} b$$

$$A\hat{x} = b = B^T Q_1^T Q_1 B^{-T} b$$

Let  $F = Q_2$ , spanning nullspace,

for  $x = Q_2 z$ ,  $Ax = 0$ !

General form,  $x = \hat{x} + Q_2 z$ !

- Apply **LU factorization**,

$$A^T = LU$$

where,  $L \in R^{n \times m}$  a unit lower triangle (applying row pivoting Alg.),  $U \in R^{m \times m}$  upper triangle.

Substituting factors,

$$\text{Let } L = \begin{pmatrix} L_1 \\ L_2 \end{pmatrix},$$

$$F = \begin{pmatrix} -L_1^{-T} L_2^T \\ I \end{pmatrix}, \quad \hat{x} = \begin{pmatrix} L_1^{-T} U b \\ 0 \end{pmatrix}$$

where  $L_1 \in R^{m \times m}$  unit lower triangle,  $L_2 \in R^{(n-m) \times m}$

- Solving ill-conditioned problem

- Some singular values of  $A$  tend to zero!

- Let  $A = U\Sigma V^T$ , Tolerable condition number,

$$k(A) = \frac{\lambda_1}{\lambda_r}, r \leq \min(m, n)$$

- Let the singular values below  $\lambda_r$  be 0!

$$A_r = U \begin{pmatrix} \Sigma_r & 0 \\ 0 & 0 \end{pmatrix} V^T$$

- Solution  $x$  of the smallest norm,

$$x = V \begin{pmatrix} \Sigma_r^{-1} & 0 \\ 0 & 0 \end{pmatrix} U^T b$$

- The Particular solution is some pseudoinverse solution!

- Iterative methods for  $Ax = b$ 
  - Select a nonsingular matrix,  $M$ , and partition  $A$  as,  $A = M - N$ , and  $M^{-1}A = I - M^{-1}N$
  - Define iterative scheme, initial guess,  $x^{(0)}$ ,
 
$$Mx^{(k+1)} = b + Nx^{(k)}, k = 0, 1, \dots$$

$$x^{(k+1)} = M^{-1}b + M^{-1}Nx^{(k)}$$
  - Iteration matrix,
  $M^{-1}N$  and vector  $M^{-1}b$  are precomputed!
  - Problem,
 

how to select  $M$ ?/or how to partition  $A$ ?

(partition methods)



## ○ Jacobian Iteration

Let  $M = \text{diag}(A) = D$  and  $N = -(L + U)$ ,

$$x^{(k+1)} = D^{-1}(b - (L + U)x^{(k)})$$

$$x_i^{(k+1)} = \frac{1}{a_{ii}} \left( b_i - \sum_{j=1}^{i-1} a_{ij}x_j^{(k)} - \sum_{j=i+1}^n a_{ij}x_j^{(k)} \right)$$

Note that, computing variables,  $x_i^{(k+1)}$ , in parallel

If  $A$  is diagonal dominant, for any  $x^{(0)}$ , it can converges!

Convergence is very slow!

○ Gauss-Seidel iteration,

To speed up convergence,

Let  $M = D + L$ , and  $N = -U$ ,

Iteration scheme,

$$(D + L)x^{(k+1)} = b - Ux^{(k)}$$
$$x^{(k+1)} = D^{-1}(b - Lx^{(k+1)} - Ux^{(k)})$$
$$x_i^{(k+1)} = \frac{1}{a_{ii}} \left( b_i - \sum_{j=1}^{i-1} a_{ij}x_j^{(k+1)} - \sum_{j=i+1}^n a_{ij}x_j^{(k)} \right)$$

If  $A$  is SPD or diagonal dominant,

It can converge for any  $x^{(0)}$ .

- Successive over-relaxation (SOR)

To further accelerate the convergence,

$$\text{Let } M = \frac{1}{w}D + L \text{ and } N = \left(\frac{1}{w} - 1\right)D - U$$

Iteration scheme,

$$\left(\frac{1}{w}D + L\right)x^{(k+1)} = \left(\left(\frac{1}{w} - 1\right)D - U\right)x^{(k)} + b$$

$$x^{(k+1)} = (1 - w)x^{(k)} + w\left(D^{-1}(b - Lx^{(k+1)} - Ux^{(k)})\right)$$

The 2<sup>nd</sup> term is the Gauss-Seidel iteration!

$$x_i^{(k+1)} = (1 - w)x_i^{(k)} + \frac{w}{a_{ii}} \left( b_i - \sum_{j=1}^{i-1} a_{ij}x_j^{(k+1)} - \sum_{j=i+1}^n a_{ij}x_j^{(k)} \right)$$

$$0 < w < 2!$$

Appropriate  $w$  will effectively accelerate the convergence!

- Summary

- Gaussian Elimination
- LU factorization (pivoting, Cholesky factorization)
- Iterative methods (Jacobi's method, Gaussian-Seidel, SOR)

# 2. Linear Least Squares Methods

- Least squares problem
  - Standard approach to the approximate solution of overdetermined systems;
  - Application: Data fitting!
  - Linear least squares problems,

$$\sum_{j=1}^n x_{ij}\beta_j = y_i, i = 1, \dots, m, m > n$$

Or, matrix form,  $X\beta = Y, X \in R^{m \times n}, Y \in R^n,$

Quadratic optimization form,

$$\beta^* = \arg \min_{\beta} F(\beta) = \|Y - X\beta\|^2$$

– Normal equation,

$$\frac{\partial F}{\partial \beta} = 0$$
$$(X^T X)\beta = X^T Y$$

where,  $X$  is full rank!

There exists Closed-Form Solution!

– Another method for full rank  $X$ ,

$$\min_{\beta} \|Y - X\beta\|^2$$
$$= \|UDV^T\beta - Y\|^2 = \|DV^T\beta - U^TY\|^2$$

Let  $\beta' = V^T\beta$  and  $Y' = U^TY$ ,

$$\beta' = \operatorname{argmin} \|D\beta' - Y'\|$$
$$\beta = V\beta'$$

- When  $X$  is not full rank,

$$\begin{aligned} X &= (U_1, U_2) \begin{pmatrix} D \\ 0 \end{pmatrix} V^T \\ \min_{\beta} \|Y - X\beta\|^2 & \\ &= \|U^T(Y - X\beta)\|^2 \\ &= \|DV^T\beta - U_1^T Y\|^2 + \|U_2^T Y\|^2 \end{aligned}$$

- Recall the solution form of full rank  $X$ ,

$$\begin{cases} \beta^* = VD^{-1}U_1^T Y \\ \beta^* = \sum_{i=1}^n \frac{u_i^T Y v_i}{\sigma_i} \end{cases}$$

– When  $\sigma_{i,\dots,n} = 0$ , general solution,

$$\beta = \sum_{\sigma_i \neq 0} \frac{u_i^T Y v_i}{\sigma_i} + \sum_{\sigma_i = 0} \alpha_i u_i$$

where, arbitrary factors  $\alpha_i$ .

○ The particular solution with the smallest norm,

$$\beta^* = \sum_{\sigma_i \neq 0} \frac{u_i^T Y v_i}{\sigma_i}$$



– Another method for **rank-deficient system (Pseudo-inverse)**

For a deficient rank system  $Ax = b$ , let  $A = UDV^T$ . Define pseudo-inverse of  $D$  as,

$$d_{ii}^+ = \begin{cases} 0, & \text{if } d_{ii} = 0 \\ d_{ii}^{-1}, & \text{otherwise} \end{cases}$$

○ The pseudo-inverse of  $A$ ,  $A^+ = VD^+U^T$ , i.e.

$$x = A^+b$$

– **Pseudoinverse of a full rank system  $Ax = b$**

○ Case 1,  $A \in R^{m \times n}$ ,  $m > n$ ,  $\text{Rank}(A) = n$

$$A^+ = (A^T A)^{-1} A^T$$

○ Case 2,  $n > m$ ,  $\text{Rank}(A) = m$

$$A^+ = A^T (AA^T)^{-1}$$

$$A \left( A^T (AA^T)^{-1} b \right) = b \rightarrow x = A^T (AA^T)^{-1} b$$

- Weighted Least Squares

- Add a diagonal matrix  $W$ ,

$$\begin{aligned}\beta^* &= \mathop{\text{arg min}}_{\beta} S(\beta) = (Y - X\beta)^T W (Y - X\beta) \\ &= \sum_{i=1}^n w_i (y_i - X_i^T \beta)^2\end{aligned}$$

- Normal equation,

$$(X^T W X) \beta = X^T W Y$$

- Solution,

$$\beta^* = (X^T W X)^{-1} X^T W Y$$

– Define weight  $W$ ,

○ Example of fitting scattered points,

$$y_i = \beta_0 + \beta_1 x_i + \cdots + \beta_m x_i^m + \varepsilon_i$$

Define  $X$  as a polynomial basis vector,

$$\begin{bmatrix} y_1 \\ y_2 \\ y_3 \\ \vdots \\ y_n \end{bmatrix} = \begin{bmatrix} 1 & x_1 & x_1^2 & \cdots & x_1^m \\ 1 & x_2 & x_2^2 & \cdots & x_2^m \\ 1 & x_3 & x_3^2 & \cdots & x_3^m \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & x_n & x_n^2 & \cdots & x_n^m \end{bmatrix} \begin{bmatrix} \beta_0 \\ \beta_1 \\ \beta_2 \\ \vdots \\ \beta_m \end{bmatrix} + \begin{bmatrix} \varepsilon_1 \\ \varepsilon_2 \\ \varepsilon_3 \\ \vdots \\ \varepsilon_n \end{bmatrix},$$

Define  $W$  by a Gaussian,  $W = \text{diag} \left( e^{-\frac{\|x_i - \bar{x}\|^2}{h^2}} \right)$   
(polynomial regression!)

- Moving Least Squares

- Basic idea, each point  $p$  is fitted by weighted Least squares!
- Notations,

$X$  is defined by Polynomial base  $P(p) = (1, x, y, \dots)^T$ ,

e.g.  $P^*(p_i) = (1, x_i, x_i^2, \dots)^T, i = 1, \dots, n, p_i = (x_i) \quad (1D),$

$$P^*(p_i) = (1, x_i, y_i, x_i^2, x_i y_i, y_i^2, \dots)^T, p_i = (x_i, y_i) \quad (2D),$$

$$P^*(p_i) = (1, x_i, y_i, z_i, x_i^2, x_i y_i, x_i z_i, y_i^2, y_i z_i, z_i^2, \dots)^T, p_i = (x_i, y_i, z_i) \quad (3D),$$

Let weight at point  $p$  be  $W(p - p_i) = \text{diag}(w_i(p - p_i))$ , and  $w_i(\cdot)$  may be a gaussian!

Let  $p_i$  be particular points (/control points).

- Step 1: Polynomial approximation if  $\beta$  is given,

$$f(p) = P(p)\beta$$

- Step 2: Determine Surface if  $\beta$  is unknown,

$$f(p_i) = P^*(p_i)\beta, i = 1, \dots, n$$

e.g. implicit function  $f(p)$  and vector  $\beta$

– Moving least squares problem,

Define the functional  $F$  on a point  $p$ ,

$$\min_{\beta} F = \frac{1}{2} \sum_{i=1}^n w_i (p - p_i) (P^*(p_i)\beta - f_i)^2$$

Matrix form,

$$F = \frac{1}{2} (P^*\beta - f^*)_{1 \times n}^T W (p - p_i) (P^*\beta - f^*)_{n \times 1}, i = 1, \dots, n$$

$n$  control points!

○ 1<sup>st</sup> order derivative,

$$P^{*T} W P^* \beta = P^{*T} W f^*$$

○ The local solution on  $p$ ,

$$\beta^* = (P^{*T} W P^*)^{-1} P^{*T} W f^*$$

## – Implementation

- Step 1: Approximation solution on some point  $p$ ,

$$f(p) = P(p)\beta^*$$

$$= P(p)(P^{*T}W(p - p_i)P^*)^{-1}P^{*T}W(p - p_i)f^*$$

$$i = 1 \dots n$$

Note that  $P(p)$  is different from  $P^{*T}$ , since  $P^{*T}$  only contains the control points  $p_i$ . Similarly  $f(p)$  is different from  $f^*$ !

- Step 2: MLS Shape function on any point  $p$ ,

$$\varphi(p) = P(p)(P^{*T}W(p - p_i)P^*)^{-1}P^{*T}W(p - p_i)f^*$$

$$i = 1 \dots n$$

Note that, (1)  $\varphi(p)$  are unknown in closed form! (2)  $W$  of  $\varphi(p)$  depends on  $(p - p_i), i = 1 \dots n$ !

Therefore, for a new point  $p$ , we have to re-compute  $\varphi(p)$  accordingly, i.e. updating  $\beta^*$  for every new point  $p$ !

# 3. Constrained Least Squares

- Least squares solution of **homogeneous Eqs.**

- Find  $x$  minimizing  $\|Ax\|$  subject to  $\|x\| = 1$

- Take SVD,  $\|UDV^T x\| = \|DV^T x\|$

and

$$\|x\| = \|V^T x\|$$

Thus,  $\min_x \|Ax\|$  is equivalent to,

$$\min_y \|Dy\|, \quad s.t. \|y\| = 1, \quad y = V^T x$$

There exists a solution  $y = (0, \dots, 0, 1)^T$ , which implies that  $x = Vy$  and the last column of  $V$  is the least squares solution of  $Ax = 0$ . (it can be proven by the eigenvalue decomposition of  $A^T A$ )

- Solving  $\min_x \|Ax\|^2$ 
  - Recall Rayleigh quotient,  $\min_x \frac{x^T A^T A x}{x^T x}$ , if  $\|x\| \neq 1$ ,
  - SVD of A,  $A = U\Sigma V^T$ ,  $\min_x x^T V \Sigma^2 V^T x$
  - When the smallest eigenvalue is zero, it reaches the minimal at x as the eigenvector corresponding to the smallest eigenvalue.



- Remark

- This is a standard form, which most constrained least square problems can be converted into!
- The basic idea is to solve a few of typical least square problems as standard forms and then convert many least square problems into some appropriate standard forms through matrix techniques!

- Inequality constraints,

- Solve  $x = \operatorname{argmin} \|Ax - b\|^2$ , s. t.  $\|x\| \leq a$

- Take SVD to  $A$ ,  $A = U\Sigma V^T$ ,

$$\|U\Sigma V^T x - b\| = \|\Sigma \hat{x} - \hat{b}\|$$

where,  $\hat{x} = V^T x$  and  $\hat{b} = U^T b$ ,  $r = \operatorname{rank}(A)$

- If  $\sum_{i=1}^r \left(\frac{\hat{b}_i}{\sigma_i}\right)^2 > a^2$ ,

find  $\rho^*$  such that  $\sum_{i=1}^r \left(\frac{\sigma_i \hat{b}_i}{\sigma_i^2 + \rho^*}\right)^2 = a^2$ ,

update,  $x = (v_i)_{i=1:r} \left(\frac{\sigma_i \hat{b}_i}{\sigma_i^2 + \rho^*}\right)_{i=1:r}^T$

- If  $\sum_{i=1}^r \left(\frac{\hat{b}_i}{\sigma_i}\right)^2 \leq a^2$ ,

update  $x = (v_i)_{i=1:r} \left(\frac{\hat{b}_i}{\sigma_i}\right)_{i=1:r}^T$

– Find  $x$ ,  $\min_x \|Ax - b\|$  subject to  $\|Bx - d\| \leq a$

○ Let  $A \in R^{n_1 \times m}$ ,  $B \in R^{n_2 \times m}$ , take Generalized SVD to  $\begin{pmatrix} A \\ B \end{pmatrix}$

Convert to,

$$\min_y \|D_A y - \hat{b}\|, \text{ s. t. } \|D_B y - \hat{d}\| \leq a$$

$$\text{where } y = X^{-1}x, \hat{b} = U_1^T b, \hat{d} = U_2^T d$$

○ Apply Lagrange multipliers,

$$\min_{\lambda, y} \|D_A y - \hat{b}\|^2 + \lambda (\|D_B y - \hat{d}\|^2 - a^2)$$

○ Take gradient on it w.r.t.  $y$ ,

$$(D_A^T D_A + \lambda D_B^T D_B)y = D_A^T \hat{b} + \lambda D_B^T \hat{d}$$

○ Solution,

$$y_i(\lambda) = \begin{cases} \frac{\alpha_i \hat{b}_i + \lambda \beta_i \hat{d}_i}{\alpha_i^2 + \lambda \beta_i^2}, & i = 1 \dots p \\ \frac{\hat{b}_i}{\alpha_i}, & i = p + 1, \dots m \end{cases}$$

Lagrange parameter  $\lambda$ ,

$$\lambda = \operatorname{argmin} \|D_B y(\lambda) - \hat{d}\|^2 - a^2$$

Note that it is 1-dimensional search!

- Equality Constraints

- Find  $x$  minimizing  $\|Ax\|$  subject to  $\|x\| = 1$  and  $Cx = 0$ .

- Assume  $\text{rank}(C) = r$ ,  $C = UDV^T$

- Note that the row space of  $C$  is generated by the first  $r$  rows of  $V^T$ .

- Orthogonal complement of row space,  $C^\perp$ , is the null space of  $C$ , i.e. the last columns of  $V$  corresponding to 0-eigenvalue

$$CC^\perp = 0$$

For  $x$  of constraint  $Cx = 0$ , (involving constraint into cost Fun!)

$$x = C^\perp y \text{ and } \|x\| = \|C^\perp y\| = \|y\|$$

- Convert to,  $\min_y \|AC^\perp y\|$  subject to  $\|y\| = 1$

(comparing it with a standard form!)

– Find  $x$  minimizing  $\|Ax\|$  subject to  $\|x\| = 1$  and  $x = G\hat{x}$  ( $x$  and  $\hat{x}$  are unknown)

- Assume  $\text{rank}(G) = r$ ,  $G = UDV^T$
- Note that the row space of  $G^T$  is generated by the first  $r$  rows of  $U^T$ .
- Orthogonal complement of row space,  $G^\perp$ , is *the null space of  $G^T$* , i.e.

$$G^T G^\perp = 0$$

- For  $x$  of  $G^\perp x = 0$ ,  $x$  is the solution of  $x = G\hat{x}$ , i.e. nullspace! Only consider nullspace here!
- Thus convert constraint  $x = G\hat{x}$  to  $G^\perp x = 0$ , i.e.

$$\min_x \|Ax\|, \quad \text{s. t. } G^\perp x = 0 \text{ and } \|x\| = 1$$

(this is the previous problem!)

- If  $\hat{x}$  is required too,  $\hat{x} = G^+ x$ .

– Find  $x$ ,  $\min_x \|Ax - b\|$ , s. t.  $Bx = d$

where  $B \in R^{p \times n}$ ,  $A \in R^{m \times n}$ ,  $d \in R^p$ ,  $\text{rank}(B) = p$ ,  $A$  and  $B$  full rank;

○ Apply QR factorization to  $B$ ,  $Q^T B^T = \begin{pmatrix} T \\ 0 \end{pmatrix}$ ,  $Q \in R^{n \times n}$ ,  $T \in R^{p \times p}$

Let  $AQ = (A_1, A_2)$ ,  $A_1 \in R^{m \times p}$ ,  $A_2 \in R^{m \times (n-p)}$ ,  $Q^T x = \begin{pmatrix} y \\ z \end{pmatrix}$ ,  $y \in R^p$ ,  $z \in R^{(n-p)}$ ;

○ Then,

$$Ax = A_1 y + A_2 z \text{ and } T^T y = d$$

○ Solution,

To solve  $y$  by backward substitution.

$$z = \arg \min_z \|A_2 z - (b - A_1 y)\|^2$$

$$x = Q \begin{pmatrix} y \\ z \end{pmatrix}$$

– Find  $x$ , minimizing  $\|Ax\|$  subject to  $\|Cx\| = 1$

- Let  $C = UDV^T$ ,  $\|Cx\| = 1$  is equivalent to,

$$\|DV^T x\| = 1$$

Convert to,  $\min_y \|AVy\|$  s. t.  $\|Dy\| = 1, y = V^T x$

- Let  $\text{rank}(D) = r$ , and  $A' = A(V_{:,1:r}, V_{:,r+1:n}) = (A'_1, A'_2), D = \begin{pmatrix} D_{1:r,:} \\ 0 \end{pmatrix}, y = (V_{:,1:r}, V_{:,r+1:n})^T x = \begin{pmatrix} y_1 \\ y_2 \end{pmatrix}$

$\|A'y\|$  is equivalent to  $\|A'_1 y_1 + A'_2 y_2\|$  partitioning  $A'$  based on rank, i.e.  $A'y = (A'_1, A'_2) \begin{pmatrix} y_1 \\ y_2 \end{pmatrix}$

The constraint is changed as,  $\|D_{1:r} y_1\| = 1$

- Substituting  $y_2 = -A_2'^+ A_1' y_1$  to  $\|A_1' y_1 + A_2' y_2\|$ ,  
 $\min_{x'} \|(I - A_2' A_2'^+) A_1' D_1^{-1} x'\|, \text{ s. t. } \|x'\| = 1$

where,  $x' = D_1 y_1$ , (the previous problem!)



# 4. Nonlinear Least Squares Methods

- General least squares problems,
  - Residual form (real function),

$$\min_{x \in \mathbb{R}^n} F(x) = \frac{1}{2} \sum_{i=1}^m r_i^2(x), m \geq n$$

where  $r_i(x(t_i)) = y_i - f(x(t_i))$  (e.g. data fitting). Or,

$$\min_{x \in \mathbb{R}^n} \frac{1}{2} \|r(x)\|_2^2, r \in \mathbb{R}^m$$

- 1<sup>st</sup> order derivative,

$$\nabla F(x) = \sum_{i=1}^m r_i(x) \nabla r_i(x) = J(x)^T r(x)$$

where, Jacobian  $J(x) = \begin{pmatrix} \frac{\partial r_i}{\partial x_j} \end{pmatrix}_{\substack{i=1, \dots, m \\ j=1, \dots, n}}$

- Apply 1<sup>st</sup> order Taylor expansion to  $r(x)$  around optimal solution  $x^*$ ,

$$\min_x \|r(x^*) + J(x^*)(x - x^*)\|_2$$

This is to solve linear system,

$$\min_x \|J^*x - b^*\|^2$$

Computing pseudoinverse requires  $J^{*T}J^*$  is full rank!

- Assumption of iteration scheme,

The sequence of Jacobian  $\{J_k\}$  is always full rank!

Rising Question,  $J_k$  may be rank deficient!

– Hessian,

$$\begin{aligned}\nabla^2 F(x) &= \sum_{j=1}^m \nabla r_j(x) \nabla r_j(x)^T + \sum_{j=1}^m r_j(x) \nabla^2 r_j(x) \\ &= J(x)^T J(x) + \sum_{j=1}^m r_j(x) \nabla^2 r_j(x)\end{aligned}$$

Note that  $\nabla J(x) = \nabla^2 r(x)$

The 1<sup>st</sup> part dominates the others in Hessian;

The 2<sup>nd</sup> part is highly nonlinear (**usually cancelled**)!

– Challenging issue is how to approximate Hessian!

- Gaussian-Newton method

- Recall Newton method

- Jacobin  $J(x)$  is derivative of residual error,

$$J(x)_{ij} = \frac{\partial r_i(x)}{\partial x_j}$$

- Gradient is derivative of  $F(x)$ ,

$$\nabla F(x)_j = \frac{\partial f(x)}{\partial x_j} = \sum_{i=1}^n r_i(x) \frac{\partial r_i(x)}{\partial x_j}$$

- 1<sup>st</sup> order derivative (necessary condition)

$$\nabla F(x) = J(x)^T r(x) = 0$$

- If  $F(x)$  is  $C^2$  continuous, Newton's Method updates  $x$  by,

$$\begin{aligned}x^{(k+1)} &= x^{(k)} - \left( \nabla^2 F(x^{(k)}) \right)^{-1} \nabla F(x^{(k)}) \\ &= x^{(k)} - \left( J_k^T J_k + H_k \right)^{-1} J_k^T r_k\end{aligned}$$

Searching direction  $p^{(k)} = x^{(k+1)} - x^{(k)}$ , satisfying,

$$(J_k^T J_k + H_k)p^{(k)} = -J_k^T r_k$$

$H_k$  is the 2<sup>nd</sup> term of Hessian;

- If  $H_k$  is neglected,

**This is required by Gauss-Newton method!**

- The kernel of GN is its Normal Equation,

$$J_k^T J_k p^{(k)} = -J_k^T r_k$$

which yields **the GN iterative scheme** as,

$$x^{(k+1)} = x^{(k)} - (J_k^T J_k)^{-1} J_k^T r_k$$

- When  $J_k$  is full rank, Normal Equation is regarded as a linear LS problem, i.e. “ $Ax = b$ ”!

$$\min_{p^{(k)}} \|J_k p^{(k)} + r_k\|^2$$

This means that it is unnecessary to compute Hessian approximation  $J_k^T J_k$ . (There exist QR/SVD algorithms for this linear LS problem!)

**Is it true a linear LS solution is a stable solution?** (note that GN ignores the 2<sup>nd</sup> term of Hessian! We need to evaluate the influence of the 2<sup>nd</sup> term!)

– Question: GN iterative scheme convergences?

- This is a descent step with direction  $p^{(k)}$ . Consider 1<sup>st</sup> order Taylor expansion,

$$F(x) = F(x^{(k)}) + \nabla F(x^{(k)})^T p^{(k)}$$

which is expected to  $F(x) - F(x^{(k)}) \leq 0$ . So,

$$\begin{aligned} p^{(k)T} \nabla F(x^{(k)}) &= p^{(k)T} J_k^T r_k = -p^{(k)T} J_k^T J_k p^{(k)} \\ &= -\|J_k p^{(k)}\|^2 \leq 0 \end{aligned}$$

This means, reasonable solution  $x^{(k)}$  leads to

$$J_k p^{(k)} = 0$$

further leads to  $J_k^T r_k = \nabla F(x^{(k)}) = 0$ ;

○ Convergence,

$$\begin{aligned} & x^{(k)} + p^{(k)} - x^* \\ &= x^{(k)} - x^* - (J^T J)^{-1} \nabla F(x^{(k)}) \\ &= (J^T J)^{-1} \left( J^T J(x^{(k)}) (x^{(k)} - x^*) - \left( \nabla F(x^{(k)}) - \nabla F(x^*) \right) \right) \end{aligned}$$

Note  $\nabla F(x^{(k)}) = J_k^T r_k$ ,  $\nabla F(x^*) = 0$ .

Recall the definition of Hessian Matrix,

$$\begin{aligned} & \nabla F(x^{(k)}) - \nabla F(x^*) \\ &= \int_0^1 J \left( x^* + t(x^{(k)} - x^*) \right)^T J \left( x^* + t(x^{(k)} - x^*) \right) (x^{(k)} - x^*) dt \\ & \quad + \int_0^1 H \left( x^* + t(x^{(k)} - x^*) \right) (x^{(k)} - x^*) dt \end{aligned}$$

Submitting it into iterative scheme yields.....



- Norm form,

$$\begin{aligned} & \|x^{(k)} + p^{(k)} - x^*\| \\ & \leq \int_0^1 \left\| (J^T J)^{-1} H \left( x^* + t(x^{(k)} - x^*) \right) \right\| \|x^{(k)} - x^*\| dt + O \left( \|x^{(k)} - x^*\|^2 \right) \\ & \approx \left\| (J^T J(x^*))^{-1} H(x^*) \right\| \|x^{(k)} - x^*\| + O \left( \|x^{(k)} - x^*\|^2 \right) \end{aligned}$$

It is desirable,  $\left\| (J^T J(x^*))^{-1} H(x^*) \right\| < 1!!!$

(i.e. if satisfying it, GN works well; otherwise to be failed! This is the assumption of GN!)

In practice, we don't know  $H(x^*)!$

Recall the linear system,  $\min_{p^{(k)}} \|J_k p^{(k)} + r_k\|^2$

We have to determine in WHICH scenario the linear LS solution is stable!!!

– To solve Linear LS problem,

$$\min_{p^{(k)}} \|J_k p^{(k)} + r_k\|^2$$

○ Case 1: overdetermined problem

Let Jacobian  $J \in R^{m \times n}$ , direction  $p \in R^n$ , and  $m > n$ ;  
and  $x^*$  be a minimizer, to solve  $p$  by,

$$J_k p^{(k)} = r_*$$

i.e.  $p^{(k)} = x^* - x^{(k)}$ ,  $r_*$  residual error w.r.t. optimal  
solution:  $y - y^*$ !

➤ Consider  $J_k^T J_k p^{(k)} = J_k^T r_*$

Change it as,

$$J_k^T J_k p^{(k)} = (J_k^T - J_*^T) r_*$$

Because  $J_*^T r_* = 0$ , (this is the 1<sup>st</sup> order necessary condition!)

$$\begin{aligned} J_k^T r_* &= (J_* + \nabla J_*(x - x^*))^T r_* + O(\|\varepsilon\|^2) \\ &= (x - x^*)^T \nabla J_*^T r_* + O(\|\varepsilon\|^2) \end{aligned}$$

Recall,

$$\nabla^2 F(x) = J(x)^T J(x) + \sum_{j=1}^m r_j(x) \nabla^2 r_j(x)$$

We have,

$$\nabla J_*^T r_* = \nabla^2 F(x^*) - J_*^T J_*$$

Further lead to

$$J_k^T r_* = (x - x^*)^T (\nabla^2 F(x^*) - J_*^T J_*) + O(\|\varepsilon\|^2)$$

➤ Conclude the estimation upper bound

(1) If  $\|r_*\|$  is large,

$$\|J_k^T J_k p^{(k)}\| = \|J_k^T r_*\| \leq \|(\nabla^2 F(x^*) - J_*^T J_*)\| \|r_*\| + O(\|\varepsilon\|^2)$$

This says that even for a large residual problem, it can converge fast if it is not very **nonlinear**!

(2) If  $\|x - x^*\|$  is large (i.e. distant point  $x$ )

$$\begin{aligned} \|J_k^T J_k p^{(k)}\| &= \|J_k^T r_*\| \\ &\leq \|(\nabla^2 F(x^*) - J_*^T J_*)\| \|x - x^*\| + O(\|\varepsilon\|^2) \end{aligned}$$

This says that even for initial  $x$  far from  $x^*$  or large residual problem (but very **nonlinear**), there is no reason to expect GN to work well because iteration cannot decrease the upper bound!

○ Case 2: Underdetermined problem ( $m < n$ )

➤ Applying SVD to  $J_k$  yields pseudoinverse  $J_k^+$ ,

$$p^{(k)} = J_k^+ r_k$$

The pseudoinverse solution is a particular solution but not an unique solution,

The general solution is,

$$p = J_k^+ r_k + Fz$$

i.e. there is no unique solution.

Note that Gauss-Newton cannot guarantee to converge to an unique solution since it is underdetermined!

- Damped Gauss-Newton

- Recall Armijo Rule

Sufficient decrease: for  $F(x^k)$ , actual decrease is equal to the predicated decrease up to a scale.

Search on a line from  $x^k$  in direction of locally decreasing  $F(x)$ ,

Start with  $m = 0$  then increase  $m$  until sufficient decrease is achieved, i.e.  $\lambda = \alpha^m, \alpha \in (0,1)$

This is called “backtracking”/”pullbacks”,

For each  $m$ , a new function evaluation is required!

- Recall Gauss-Newton iterative scheme, if  $J_k$  is full rank, the search direction  $p^{(k)}$ ,

$$p^{(k)} = x^{k+1} - x^k = -(J_k^T J_k)^{-1} J_k^T r_k$$

The search direction is a descend direction!

Recall Newton Method for comparison,

$$x = x^k - \left( \nabla^2 F(x^k) \right)^{-1} \nabla F(x^k)$$

- Applying Armijo Rule to GN yields **Damped GN**

$$p^{(k)} \leftarrow \alpha^m p^{(k)}, \alpha \in (0,1)$$

For each  $m$ , the sufficient decrease is required, which results in **the iterative scheme** as,

$$x^{(k+1)} = x^{(k)} - \alpha^m (J_k^T J_k)^{-1} J_k^T r_k$$

## – Example

$i$	1	2	3	4	5	6	7
{x}	0.038	0.194	0.425	0.626	1.253	2.500	3.740
{y}	0.050	0.127	0.094	0.2122	0.2729	0.2665	0.3317

$$r_i = y_i - \frac{\beta_1 x_i}{\beta_2 + x_i} \quad i = 1, \dots, 7$$

Initial guess of variables  $\beta_1 = 0.9$  and  $\beta_2 = 0.2$ ,

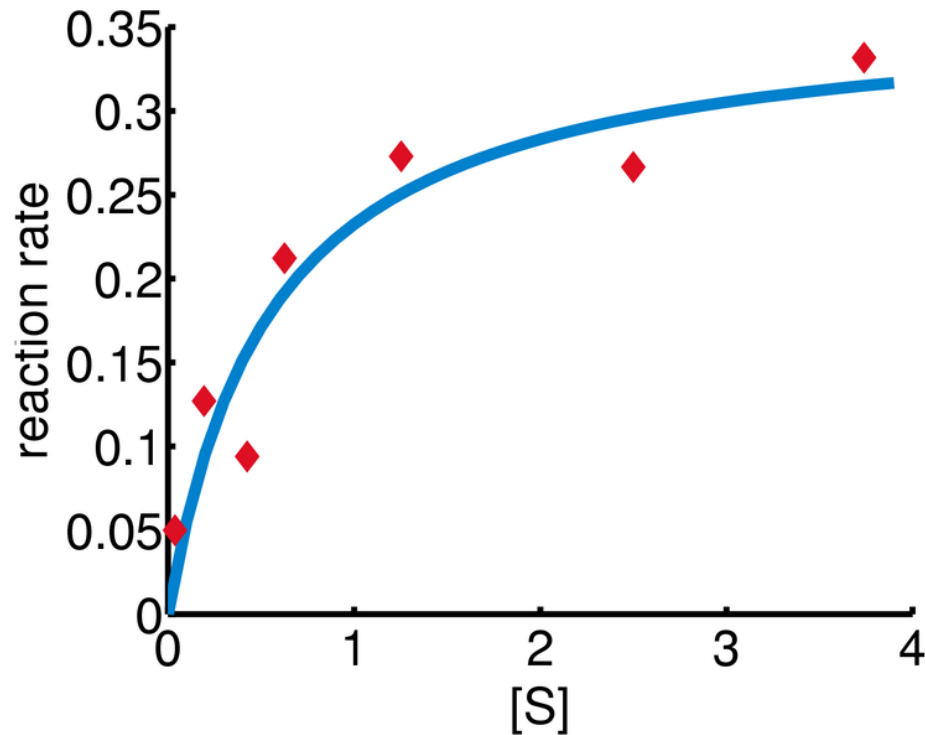
$$\text{1st order derivative, } J(x)_{ij} = \left( \frac{\partial r_i}{\partial \beta_j} \right)_{7 \times 2} =$$

$$\left( -\frac{x_i}{\beta_2 + x_i}, \frac{\beta_1 x_i}{(\beta_2 + x_i)^2} \right)$$



– Plot result

5 iterations, convergence to  $\beta^* = (0.362, 0.556)$



- Levenberg-Marquardt method

- Jacobian,  $J(x)$ , may be not full rank!

- The same idea, remove the 2<sup>nd</sup> order term of Hessian;  
Recall Gauss-Newton's method, compute the searching direction  $p$ ,

$$J_k^T J_k p^{(k)} = -J_k^T r_k$$

- **Basis idea: within a small region, Hessian is SPD!**

$$\min_p \frac{1}{2} \|J_k p + r_k\|^2$$

- where,  $\|p\| \leq \gamma$ , **i.e. strategy changing linear search to trust region!!!**

– Add constraints to nonlinear LS problem,

$$\min_p \frac{1}{2} \|J_k p + r_k\|^2 + \alpha_k \|p\|^2$$

where  $\alpha_k$  is Lagrange multiplier for constraint.

– Question,

If  $p$  is outside region  $\gamma_k$ ,  $J_k$  may be not full rank!

○ Modify the searching direction,

$$(J_k^T J_k + \rho I) p^{(k)} = -J_k^T r_k$$

in case of any ill-conditioned Jacobian.

- Formulate it as a least squares problem,

$$\min_p \frac{1}{2} \left\| \begin{pmatrix} J \\ \sqrt{\rho} I \end{pmatrix} p + \begin{pmatrix} r \\ 0 \end{pmatrix} \right\|^2$$

Then, update  $p^{(k)}$  and further update  $x^{(k)}$  .....

Note: bound constraint is changed to parameter  $\rho$ !

- Rising Questions,

If  $\rho$  is given, direction  $p$  can be updated by...

**Basic idea is to try any possible  $\rho$  .....**

How to determine  $\rho$ ?

If direction  $p^{(k)}$  is given, how to find  $\rho$ ?

**Note that  $\rho_k$  may vary from 0 to a large step size.**

**However, initial value,  $\rho_0 = \|J_0^T J_0\|$**

➤ Fast computation of solving  $\begin{pmatrix} J \\ \sqrt{\rho}I \end{pmatrix} p = \begin{pmatrix} r \\ 0 \end{pmatrix}$ ,

If applying QR decomposition,

$$\begin{pmatrix} J \\ \sqrt{\rho}I \end{pmatrix} = Q_\rho \begin{pmatrix} U_\rho \\ 0 \end{pmatrix}$$

where,  $U_\rho^T U_\rho = J^T J + \rho I$ ;

Re-formulate Equation as,

$$Q_\rho \begin{pmatrix} U_\rho \\ 0 \end{pmatrix} p = \begin{pmatrix} r \\ 0 \end{pmatrix}$$
$$U_\rho p = Q_\rho^T r$$

A proper  $\rho$  makes  $U_\rho$  full-rank!

Thus, we need the closed forms of  $Q_\rho$  and  $U_\rho$ !

## ➤ Speeding up QR computation!

Applying QR factoring to Jacobian yields

$$J = Q \begin{pmatrix} U \\ 0 \end{pmatrix}$$
$$\begin{pmatrix} J \\ \sqrt{\rho}I \end{pmatrix} = \begin{pmatrix} Q & \\ & I \end{pmatrix} \begin{pmatrix} U \\ 0 \\ \sqrt{\rho}I \end{pmatrix}$$

Then, taking **Givens Rotation**  $\bar{Q}$ , such that,

$$\bar{Q} \begin{pmatrix} U \\ 0 \\ \sqrt{\rho}I \end{pmatrix} = \begin{pmatrix} U_\rho \\ 0 \\ 0 \end{pmatrix}$$

Define,  $Q_\rho = \begin{pmatrix} Q & \\ & I \end{pmatrix} \bar{Q}^T$ !

Each updating  $Q_\rho$  for new attempt  $\rho$  only needs to re-compute  $\bar{Q}$ !

$U_\rho$  has been available when computing  $\bar{Q}$ !

➤ Compute direction  $p$  by  $U_\rho p = Q_\rho^T r$

- Summary

- Newton method

$$F(x) = F(x^k) + \nabla F(x^k)^T (x - x^k) + \frac{1}{2} (x - x^k)^T \nabla^2 F(x^k) (x - x^k)$$

- Gauss-Newton method,

$$F(x) = F(x^k) + \nabla F(x^k)^T (x - x^k) + \frac{1}{2} (x - x^k)^T J_k^T J_k (x - x^k)$$

- Damped Gauss-Newton method,

$$F(x) = F(x^k) + \nabla F(x^k)^T (x - x^k) + \frac{1}{2\alpha^m} (x - x^k)^T J_k^T J_k (x - x^k)$$

- Levenberg-Marquardt Alg.

$$F(x) = F(x^k) + \nabla F(x^k)^T (x - x^k) + \frac{1}{2} (x - x^k)^T (J_k^T J_k + \rho I) (x - x^k)$$

# 5. Orthogonal Distance Regression

- Total least squares problem
  - Consider a general noisy system,

$$\min_{x,E} \|(A + E)x - b\|_2^2$$

where,  $E$  and  $r = (A + E)x - b$  both are perturbations!

Comparing it with an usual noisy system,

$$\min_x \|Ax - b\|_2^2$$

**Basic idea** is to consider noises both from observation and from system!

- For nonlinear systems,

$$r_i = \varphi(x, t_i + \delta_i) - b_i, \quad i = 1, \dots, m$$

Note that there are 2 kinds of variables here,  $x, \delta$  while  $r$  is an implicit variable.



- Orthogonal distance regression,
  - Basic idea, minimizing two norms (residual  $r$  and input perturbation  $\delta$ !)

$$\min_{x, \delta_i, r_i} \frac{1}{2} \sum_{i=1}^m w_i^2 r_i^2 + d_i^2 \delta_i^2$$

where  $w_i, d_i$  are given.

Rewrite it as, (eliminating residual-implicit variable  $r$ !)

$$\min_{x, \delta_i} \frac{1}{2} \sum_{i=1}^{2m} s_i^2(x, \delta)$$

where,

$$s_i(x, \delta) = \begin{cases} w_i(b_i - \varphi(x, t_i + \delta_i)), & i = 1, \dots, m \\ d_{i-m} \delta_{i-m}, & i = m + 1, \dots, 2m \end{cases}$$

– Take 1<sup>st</sup> order derivative to  $s_i$ ,

$$\frac{\partial s_i}{\partial \delta_j} = \frac{\partial(\varphi(t_i + \delta_i, x) - b_i)}{\partial \delta_j} = 0, \quad i, j = 1, \dots, m, i \neq j$$

$$\frac{\partial s_i}{\partial x_j} = 0, \quad j = 1, \dots, n, i = 1, \dots, m$$

$$\frac{\partial s_{m+i}}{\partial \delta_j} = \begin{cases} d_i, & \text{if } i = j \\ 0, & \text{otherwise} \end{cases}$$

– Partition the Jacobian as,

$$J(x, \delta) = \begin{pmatrix} A & V \\ 0 & D \end{pmatrix}$$

where,  $A$  is the partial derivatives w.r.t  $x$ ;  $V$  and  $D$  are the partial derivatives w.r.t  $\delta$  and are diagonal matrices with size  $m$  by  $m$ .

– Normal equation,

$$\begin{pmatrix} A^T A + \rho I & A^T V \\ VA & V^2 + D^2 + \rho I \end{pmatrix} \begin{pmatrix} p_x \\ p_\delta \end{pmatrix} = - \begin{pmatrix} A^T r_x \\ Vr_x + Dr_\delta \end{pmatrix}$$

where vectors  $p_{x,\delta}$  contain  $x$  and  $\delta$ !

– This is solved by standard L-M Alg.

$$\min_{p_x, p_\delta, \rho} \frac{1}{2} \left\| \begin{pmatrix} J \\ \sqrt{\rho} I \end{pmatrix} \begin{pmatrix} p_x \\ p_\delta \end{pmatrix} + \begin{pmatrix} r_x \\ r_\delta \\ 0 \end{pmatrix} \right\|^2$$

# 6. Coursework

- Reference,
  - Image deformation using moving least squares, at <http://faculty.cs.tamu.edu/schaefer/research/mls.pdf>
  - Codes, at <http://www.mathworks.co.uk/matlabcentral/fileexchange/12249-moving-least-squares>

- Image deformation,
  - Given a set of control point pairs, do a specified affine transformation on a domain according to the control point pairs;
  - For an arbitrary point  $x$ , define the affine  $F_x$  as,

$$\min_{F_x} \sum_i w_i(x) \|F_x(p_i) - q_i\|^2$$

where, control point pairs,  $(p_i, q_i)$ ;  $F_x(p_i) \rightarrow q_i$ !

Let weight  $w_i(x) = \frac{1}{\|p_i - x\|^2}$ ;

– Affine transform,

- $F_x(p) = Mp + T \rightarrow q$
- Firstly solve  $T$  by,

$$T = \arg \min_T \sum_i w_i \langle Mp_i + T - q_i, Mp_i + T - q_i \rangle$$

Yield,  $T(x) = q^*(x) - Mp^*(x)$ ,

where,  $q^* = \sum_i w_i q_i / \sum_i w_i$  and  $p^* = \sum_i w_i p_i / \sum_i w_i$

– Rewrite Affine Transform,

- $F_x(p) = M(p - p^*) + q^*$

- Then solve  $M$  by,

$$\min_M \sum_i w_i \|M(p_i - p^*) - (q_i - q^*)\|^2$$

- The 1<sup>st</sup> order derivative w.r.t.  $M$ ,

$$(p_i - p^*) w_i \left( (p_i^T - p^{*T}) M^T - (q_i^T - q^{*T}) \right) = 0, i = 1..n$$

- Local solution at point  $x$ ,

$$M = \left( \sum_i (p_i - p^*)^T w_i (p_i - p^*) \right)^{-1} \sum_j w_j (p_j - p^*)^T (q_j - q^*)$$

- Mapping  $x$  by Affine  $M$ ,

$$x' = M(x - p^*) + q^*$$

Moreover, compute the pointwise affine  $M$  for updating!

Note that, many info of  $M$  can be pre-computed!

There exist non-uniform scaling and shear...



– Consider similarity deformation,

- Constrain,  $M^T M = sI!$
- Let  $M^T = (M_1, M_2)_{2 \times 2}$ ,
- $M_2 = M_1^\perp$
- $M_1^\perp p = M_1 p^\perp$
- $p^T M^T = p^T (M_1, M_1^\perp) = \begin{pmatrix} p^T \\ p^\perp \end{pmatrix} M_1$
- $\min_{M_1} \sum_i w_i \left\| \begin{pmatrix} (p_i - p^*)^T \\ -(p_i - p^*)^\perp \end{pmatrix} M_1 - (q_i - q^*) \right\|^2$

- Solution,

$M_1$

$$= \left( \begin{pmatrix} (p_i - p^*)^T \\ -(p_i - p^*)^\perp \end{pmatrix} \begin{pmatrix} (p_i - p^*)^T \\ -(p_i - p^*)^\perp \end{pmatrix} \right)^{-1} \begin{pmatrix} (p_i - p^*)^T \\ -(p_i - p^*)^\perp \end{pmatrix}^T (q_i - q^*)$$

$M_2$

$$= \left( \begin{pmatrix} (p_i - p^*)^T \\ -(p_i - p^*)^\perp \end{pmatrix} \begin{pmatrix} (p_i - p^*)^T \\ -(p_i - p^*)^\perp \end{pmatrix} \right)^{-1} \begin{pmatrix} (p_i - p^*)^T \\ -(p_i - p^*)^\perp \end{pmatrix}^T (-q_i + q^*)^\perp$$

- Local solution at point  $x$ ,

$$M = \frac{1}{\mu} \sum_i w_i \begin{pmatrix} (p_i - p^*)^T \\ -(p_i - p^*)^\perp \end{pmatrix} ((q_i - q^*), -(q_i - q^*)^\perp)$$

$$\text{where, } \mu = \sum_i w_i (p_i - p^*)^T (p_i - p^*)$$

- Mapping  $x$  by similarity  $M$ ,

$$x' = M(x - p^*) + q^*$$

Note that, many info of  $M$  can be pre-computed, but  $M$  should be computed point to point!

## – Rigid deformation

- $\min_{s,R} \sum_i w_i \|sR(p_i - p^*) - (q_i - q^*)\|^2$

s.t.  $R^T R = I, s = 1;$

- Local solution at point  $x$ ,

- $M = \frac{1}{\mu} \sum_i w_i \begin{pmatrix} (p_i - p^*)^T \\ -(p_i - p^*)^\perp \end{pmatrix} \left( (q_i - q^*), -(q_i - q^*)^\perp \right)$

- It is the same as the solution of similarity deformation!

- Factor  $\mu$ ,

$$\mu = \sqrt{\left( \sum_i w_i (q_i - q^*)^T (p_i - p^*) \right)^2 + \left( \sum_i w_i (q_i - q^*)^T (p_i - p^*)^\perp \right)^2}$$

Note that, we cannot pre-compute it.